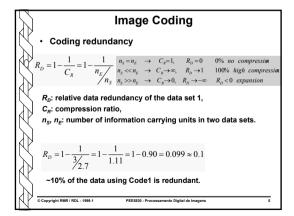
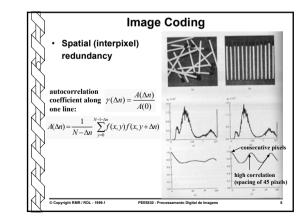
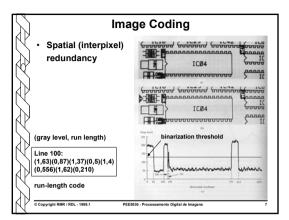
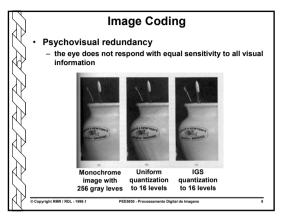


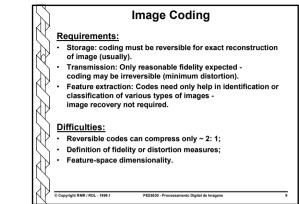
Л	Image Coding						
Y	<ul> <li>Coding redundancy - Example</li> </ul>						
D	r <sub>k</sub>	p <sub>r</sub> (r <sub>k</sub> )	Code 1	l <sub>1</sub> (r <sub>k</sub> )	Code 2	$I_2(r_k)$	$I_2(\mathbf{r}_k).\mathbf{p}_r(\mathbf{r}_k)$
Y	$r_0 = 0$	0.19	000	3	11	2	0.38
V)	$r_1 = 1/7$	0.25	001	3	01	2	0.50
12	$r_2 = 2/7$	0.21	010	3	10	2	0.42
$\sum$	$r_3 = 3/7$	0.16	011	3	001	3	0.48
$\mathbb{N}$	$r_4 = 4/7$	0.08	100	3	0001	4	0.32
Y	r <sub>5</sub> = 5/7	0.06	101	3	00001	5	0.30
N	$r_6 = 6/7$	0.03	110	3	000001	6	0.18
V.	r <sub>7</sub> = 1	0.02	111	3	000000	6	0.12
V)				L1 <sub>avg</sub> = 3			L2 <sub>avg</sub> = 2.7
	$L_{avg} = \sum_{k=0}^{L-1} l(r_k) p_r(r_k)$ Average length of the code words						
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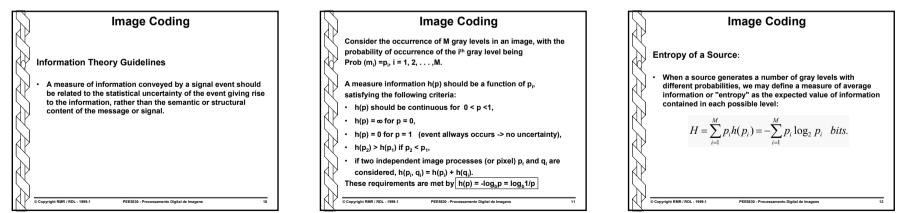


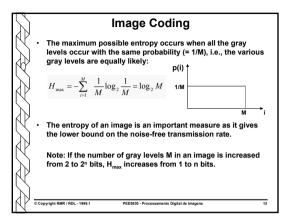


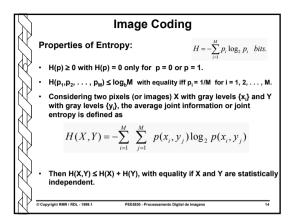


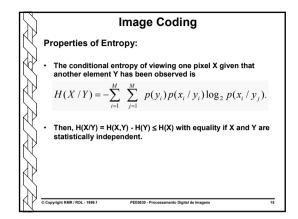


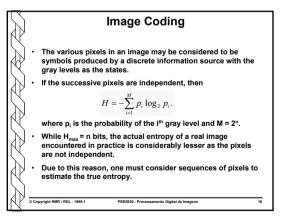


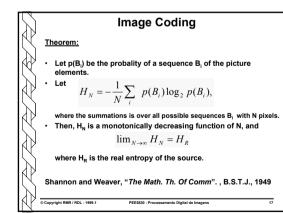


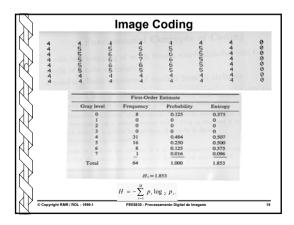


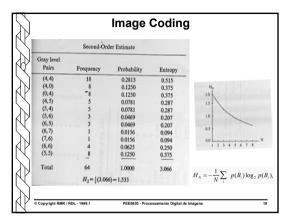


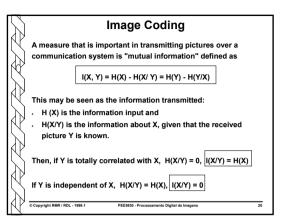


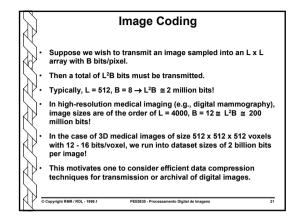












## Image Coding

### Fundamental Concepts of Coding

- A code is a mapping of words from a source alphabet into the words of a code alphabet;
- A word is a finite sequence of symbols from an alphabet.
- A code is called *distinct* if each code word is distinguishable from the other code words.
- A distinct code is *uniquely decodable* if every code word is identifiable when immersed in a sequence of code words.
- A desirable property of a uniquely decodable code is that it should be decodable on a word-to-word basis. This is ensured if no code word may be a prefix to another; the code is then instantaneously decodable.
- A code is said to be optimal if it is instantaneously decodable and has the minimum average length for a given source PDF.

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# Image Coding

### Noiseless Coding Theorem for Binary Transmission

- Given a code with an alphabet of two symbols and a source A with an alphabet of two symbols, the average length of the code words per source symbol may be made arbitrarily close to the lower bound (entropy) H(A) as desired by encoding sequences of source symbols rather than individual symbols.
- The average length L(n) of encoded n sequences is bounded

$$H_R(A) \le \frac{L(n)}{n} \le H_R(A) +$$

- Difficulty in estimating entropy lies in the fact that pixels are statistically dependent point-to-point, line-to-line, and frameto-frame.
- Computation of entropy requires the symbols to be considered in blocks over which the statistical dependence is negligible.

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#### **Image Coding** If dependencies exist over T seconds and there are F frames/seconds (considering PCM digital TV, e.g.), the block to be considered has S = T.F.B.L<sup>2</sup> a 3 x 10<sup>8</sup> bits ! (T = 5, F = 30, B = 8, L = 512). Thus joint probability functions will have to be computed with a vector length of 300 million! As this is impossible, our approaches must be limited to single pixels or small blocks, resulting in over-estimates of H(A). If the blocks of pixels are chosen so that the sequence entropy estimates converge rapidly to the limit, then blockcoding methods may provide results close to the minimum length.

 The entropy of most natural scenes; has been estimated to be less than 1 bit/pixel. Compression from 8 bits/pixel to 0.8 bits/pixel may be possible by optimal coding of blocks.

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